

\mathcal{DH}^2 -matrix compression for Helmholtz problems

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Overview

- 1 Introduction
- 2 Directional approximation
- 3 Algebraic compression
- 4 Experiments

Helmholtz equation

Goal: Solve the Helmholtz equation

$$\begin{aligned}\Delta u(x) + \kappa^2 u(x) &= 0 && \text{for all } x \in \Omega, \\ u(x) &= f(x) && \text{for all } x \in \partial\Omega.\end{aligned}$$

Challenge: Indefinite operator, pollution effect.

Approach: Boundary integral formulation

$$u(x) = \int_{\partial\Omega} g(x, y) \frac{\partial u}{\partial n}(y) dy - \int_{\partial\Omega} \frac{\partial g}{\partial n(y)}(x, y) u(y) dy \quad \text{for all } x \in \Omega$$

with fundamental solution

$$g(x, y) = \frac{\exp(i\kappa\|x - y\|)}{4\pi\|x - y\|}.$$

Challenge: Standard discretization leads to dense matrix.

Fast summation

Standard approach: Approximate g by degenerate kernel,

$$g(x, y) \approx \sum_{\nu=1}^k v_{\nu}(x) \sum_{\mu=1}^k s_{\nu\mu} \overline{w_{\mu}(y)}$$

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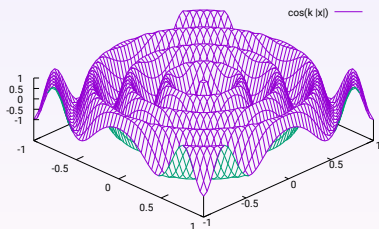
$$\int_{\partial\Omega} g(x, y) u(y) dy \approx \sum_{\nu=1}^k v_{\nu}(x) \int_{\partial\Omega} \sum_{\mu=1}^k s_{\nu\mu} \overline{w_{\mu}(y)} u(y) dy.$$

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Challenge: Helmholtz kernel oscillates rapidly if κ large.



$$g(x, y) = \frac{\exp(i\kappa\|x - y\|)}{\|x - y\|}.$$

Consequence: Standard approximation schemes require $k \gg 1$.

Experiment: \mathcal{H}^2 -matrix compression

Helmholtz matrix approximated by quasi-optimal \mathcal{H}^2 -matrix.

n	κ	k	k/κ	M/n
2048	8	33	4.1	12.5
4608	12	47	3.9	15.4
8192	16	61	3.8	18.3
18432	24	95	4.0	22.2
32768	32	132	4.1	25.4
73728	48	263	5.5	30.1
131072	64	333	5.2	36.0

Observation: Maximal rank proportional to wave number κ , storage looks like $\mathcal{O}(n \log^2 \kappa)$.

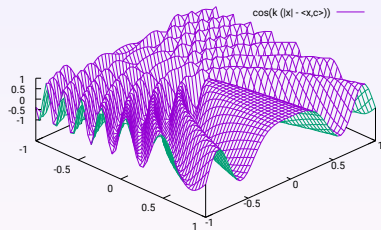
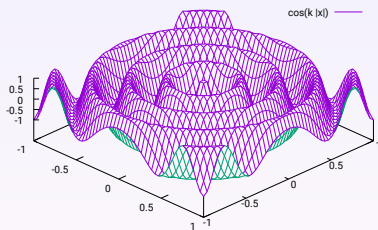
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Directional approximation

Idea: Divide spherical wave by plane wave in direction c , $\|c\| = 1$:

$$g(x, y) = \exp(i\kappa \langle x - y, c \rangle) \underbrace{\frac{\exp(i\kappa(\|x - y\| - \langle x - y, c \rangle))}{4\pi\|x - y\|}}_{=g_c(x, y)}.$$



Brandt (1991), Engquist/Ying (2007),
Messner/Schanz/Darve (2012), Bebendorf/Kuske/Venn (2015)

Directional interpolation

Approach: Apply interpolation to smoothed function g_c .

$$\tilde{g}_{\tau\sigma,c}(x, y) := \sum_{\nu,\mu=1}^k \mathcal{L}_{\tau,\nu}(x) g_c(\xi_{\tau,\nu}, \xi_{\sigma,\mu}) \mathcal{L}_{\sigma,\mu}(y).$$

Directional approximation in domain $\tau \times \sigma$ given by

$$\begin{aligned} g(x, y) &= \exp(\iota\kappa\langle x - y, c \rangle) g_c(x, y) \approx \exp(\iota\kappa\langle x - y, c \rangle) \tilde{g}_{\tau\sigma,c}(x, y) \\ &= \sum_{\nu,\mu=1}^k \underbrace{\exp(\iota\kappa\langle x, c \rangle) \mathcal{L}_{\tau,\nu}(x)}_{=: \mathcal{L}_{\tau c,\nu}(x)} g_c(\xi_{\tau,\nu}, \xi_{\sigma,\mu}) \underbrace{\exp(-\iota\kappa\langle y, c \rangle) \mathcal{L}_{\sigma,\mu}(y)}_{=: \overline{\mathcal{L}_{\sigma c,\mu}(y)}} \\ &= \sum_{\nu,\mu=1}^k \mathcal{L}_{\tau c,\nu}(x) g_c(\xi_{\tau,\nu}, \xi_{\sigma,\mu}) \overline{\mathcal{L}_{\sigma c,\mu}(y)} =: \tilde{g}_{\tau\sigma}(x, y). \end{aligned}$$

Matrix factorization

Goal: Approximate Galerkin matrix $G \in \mathbb{C}^{\mathcal{I} \times \mathcal{I}}$ given by

$$g_{ij} = \int_{\partial\Omega} \varphi_i(x) \int_{\partial\Omega} g(x, y) \varphi_j(y) dy dx.$$

Idea: If $g|_{\tau \times \sigma} \approx \tilde{g}_{\tau\sigma}$, choose

$$\hat{\tau} := \{i \in \mathcal{I} : \text{supp } \varphi_i \subseteq \tau\}, \quad \hat{\sigma} := \{j \in \mathcal{I} : \text{supp } \varphi_j \subseteq \sigma\}$$

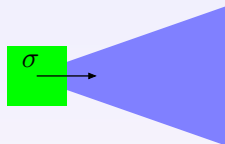
and approximate $G|_{\hat{\tau} \times \hat{\sigma}}$ by discretizing $\tilde{g}_{\tau\sigma}$ instead of g .

Result: Factorization $G|_{\hat{\tau} \times \hat{\sigma}} \approx V_{\tau\mathcal{C}} S_{\tau\sigma} V_{\sigma\mathcal{C}}^*$ with

$$v_{\tau\mathcal{C}, i\nu} = \int_{\partial\Omega} \varphi_i(x) \mathcal{L}_{\tau\mathcal{C}, \nu}(x) dx, \quad s_{\tau\sigma, \nu\mu} = g_{\mathcal{C}}(\xi_{\tau, \nu}, \xi_{\sigma, \mu}).$$

Directions

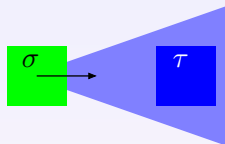
Kernel function $g_c(\cdot, y)$ for $y \in \sigma$
is smooth in a cone with axis c .



Directions

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Interactions with clusters τ inside
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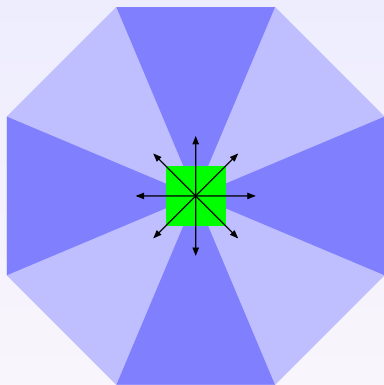


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Multiple cones have to be used to cover the entire domain.

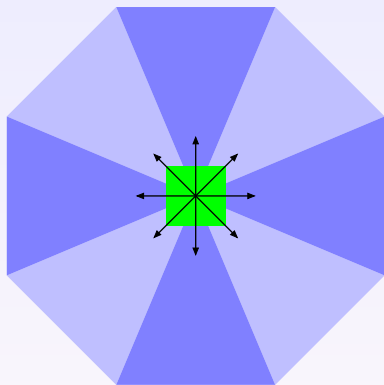


Directions

Kernel function $g_c(\cdot, y)$ for $y \in \sigma$ is smooth in a cone with axis c .

Interactions with clusters τ inside the cone can be approximated.

Multiple cones have to be used to cover the entire domain.



Approach: Choose fixed set of directions for each τ .
 $\mathcal{O}(1 + \kappa^2 \text{diam}^2(\tau))$ directions are sufficient.

Nested basis

Problem: Storing $V_{\tau c}$ for all τ, c requires $\sim n^2$ units of memory.

Solution: Organize domains τ in a **cluster tree** such that

$$\text{sons}(\tau) \neq \emptyset \quad \implies \quad \hat{\tau} = \bigcup_{\tau' \in \text{sons}(\tau)} \hat{\tau}',$$

and approximate Lagrange polynomials by weighted interpolation

$$\mathcal{L}_{\tau c, \nu} \approx \sum_{\nu'=1}^k \mathbf{e}_{\tau' \tau c, \nu' \nu} \mathcal{L}_{\tau' c', \nu'}.$$

Result: Nested representation $V_{\tau c}|_{\tau'} \approx V_{\tau' c'} E_{\tau' \tau c}$ if $\tau' \in \text{sons}(\tau)$.



Messner/Schanz/Darve (2012), B./Melenk (2015)

\mathcal{DH}^2 -matrix

Approach:

- Fix **set of directions** \mathcal{D}_τ for all clusters τ .
- Choose **admissible blocks** $(\tau, \sigma, \mathbf{c}) \in \mathcal{L}_{I \times I}^+$.
- Compute **leaf matrices** $V_{\tau\mathbf{c}} \in \mathbb{C}^{\hat{\tau} \times k}$ for all leaves τ and all $\mathbf{c} \in \mathcal{D}_\tau$.
- Find **transfer matrices** $E_{\tau'\tau\mathbf{c}} \in \mathbb{C}^{k \times k}$ for all $\tau' \in \text{sons}(\tau)$, $\mathbf{c} \in \mathcal{D}_\tau$.
- Create **coupling matrices** $S_{\tau\sigma} \in \mathbb{C}^{k \times k}$ for all $(\tau, \sigma, \mathbf{c}) \in \mathcal{L}_{I \times I}^+$.

Results:

- Storage $\mathcal{O}(kn + k^2\kappa^2 \log n)$.
- Matrix-vector multiplication in $\mathcal{O}(kn + k^2\kappa^2 \log n)$ operations.
- High-frequency case: $\kappa^2 \sim n$, complexity $\mathcal{O}(nk^2 \log n)$.



B. (2015)

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Approximation of general matrices

Given: Matrix $G \in \mathbb{C}^{\mathcal{I} \times \mathcal{I}}$ and cluster/block structure.

Goal: Find \mathcal{DH}^2 -matrix approximation of G ,
i.e., directional cluster basis ($V_{\tau\mathcal{C}}$) and coupling matrices ($S_{\tau\sigma}$).

Orthogonal basis: If we ensure $V_{\tau\mathcal{C}}^* V_{\tau\mathcal{C}} = I$, the best coupling matrices are given by orthogonal projection

$$G|_{\hat{\tau} \times \hat{\sigma}} \approx V_{\tau\mathcal{C}} \underbrace{V_{\tau\mathcal{C}}^* G|_{\hat{\tau} \times \hat{\sigma}} V_{\sigma\mathcal{C}}}_{=: S_{\tau\sigma}} V_{\sigma\mathcal{C}}^*$$

Row and column projections can be analyzed separately,

$$\begin{aligned} \|G|_{\hat{\tau} \times \hat{\sigma}} - V_{\tau\mathcal{C}} S_{\tau\sigma} V_{\sigma\mathcal{C}}^*\|_2^2 &\leq \|G|_{\hat{\tau} \times \hat{\sigma}} - V_{\tau\mathcal{C}} V_{\tau\mathcal{C}}^* G|_{\hat{\tau} \times \hat{\sigma}}\|_2^2 \\ &\quad + \|G|_{\hat{\tau} \times \hat{\sigma}} - G|_{\hat{\tau} \times \hat{\sigma}} V_{\sigma\mathcal{C}} V_{\sigma\mathcal{C}}^*\|_2^2. \end{aligned}$$

Low-rank structure

Question: What has to be approximated by $V_{\tau\mathbf{c}}$?

Answer: All blocks $G|_{\hat{\tau} \times \hat{\sigma}}$ with $(\tau, \sigma, \mathbf{c}) \in \mathcal{L}_{I \times I}^+$.

Low-rank structure

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Partial answer: All blocks $G|_{\hat{\tau} \times \hat{\sigma}}$ with $(\tau, \sigma, \mathbf{c}) \in \mathcal{L}_{\mathcal{I} \times \mathcal{I}}^+$.

Important: Directional cluster basis is **nested**.

Anything that is not approximated in the son clusters cannot be approximated in the father.

Low-rank structure

Question: What has to be approximated by $V_{\tau c}$?

Partial answer: All blocks $G|_{\hat{\tau} \times \hat{\sigma}}$ with $(\tau, \sigma, c) \in \mathcal{L}_{I \times I}^+$.

Important: Directional cluster basis is **nested**.

Anything that is not approximated in the son clusters cannot be approximated in the father.

Complete answer: All blocks $G|_{\hat{\tau} \times \hat{\sigma}^*}$ where

- $(\tau^*, \sigma^*, c^*) \in \mathcal{L}_{I \times I}^+$,
- τ is a descendant of τ^* , and
- c^* approximates c .

We collect these tuples in the set $\mathcal{R}_{\tau c} \subseteq \mathcal{L}_{I \times I}^+$.

Leaf matrices

Assume that τ is a leaf cluster and that $\mathcal{R}_{\tau c}$ is available for all $c \in \mathcal{D}_\tau$.

Goal: Find $V_{\tau c}$ with k columns such that

$$G|_{\hat{\tau} \times \hat{\sigma}^*} \approx V_{\tau c} V_{\tau c}^* G|_{\hat{\tau} \times \hat{\sigma}^*} \quad \text{for all } (\tau^*, \sigma^*, c^*) \in \mathcal{R}_{\tau c}.$$

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Reformulation: Define the submatrix

$$G_{\tau c} := G|_{\hat{\tau} \times \mathcal{F}_{\tau c}}, \quad \mathcal{F}_{\tau c} := \bigcup \{ \hat{\sigma}^* : (\tau^*, \sigma^*, c^*) \in \mathcal{R}_{\tau c} \}$$

with **farfield matrix** $G_{\tau c}$ and consider only

$$G_{\tau c} \approx V_{\tau c} V_{\tau c}^* G_{\tau c}.$$

→ Compute SVD of $G_{\tau c}$, remove small singular values.

Transfer matrices

Assume $\text{sons}(\tau) = \{\tau_1, \tau_2\}$, $\mathbf{c}_1 \in \mathcal{D}_{\tau_1}$, $\mathbf{c}_2 \in \mathcal{D}_{\tau_2}$ approximate $\mathbf{c} \in \mathcal{D}_{\tau}$.

Recursion: Ensure that $V_{\tau_1 \mathbf{c}_1}$, $V_{\tau_2 \mathbf{c}_2}$ have already been computed.

Reduced matrices given by

$$\widehat{\mathbf{G}}_{\tau \mathbf{c}} = \begin{pmatrix} V_{\tau_1 \mathbf{c}_1}^* \mathbf{G}|_{\hat{\tau}_1 \times \mathcal{F}_{\tau \mathbf{c}}} \\ V_{\tau_2 \mathbf{c}_2}^* \mathbf{G}|_{\hat{\tau}_2 \times \mathcal{F}_{\tau \mathbf{c}}} \end{pmatrix}, \quad \widehat{\mathbf{V}}_{\tau \mathbf{c}} = \begin{pmatrix} E_{\tau_1 \tau \mathbf{c}} \\ E_{\tau_2 \tau \mathbf{c}} \end{pmatrix}$$

lead to

$$\begin{aligned} \|\mathbf{G}_{\tau \mathbf{c}} - V_{\tau \mathbf{c}} V_{\tau \mathbf{c}}^* \mathbf{G}_{\tau \mathbf{c}}\|_2^2 &\leq \|\mathbf{G}|_{\hat{\tau}_1 \times \mathcal{F}_{\tau \mathbf{c}}} - V_{\tau_1 \mathbf{c}_1} V_{\tau_1 \mathbf{c}_1}^* \mathbf{G}|_{\hat{\tau}_1 \times \mathcal{F}_{\tau \mathbf{c}}}\|_2^2 \\ &\quad + \|\mathbf{G}|_{\hat{\tau}_2 \times \mathcal{F}_{\tau \mathbf{c}}} - V_{\tau_2 \mathbf{c}_2} V_{\tau_2 \mathbf{c}_2}^* \mathbf{G}|_{\hat{\tau}_2 \times \mathcal{F}_{\tau \mathbf{c}}}\|_2^2 \\ &\quad + \|\widehat{\mathbf{G}}_{\tau \mathbf{c}} - \widehat{\mathbf{V}}_{\tau \mathbf{c}} \widehat{\mathbf{V}}_{\tau \mathbf{c}}^* \widehat{\mathbf{G}}_{\tau \mathbf{c}}\|_2^2. \end{aligned}$$

The first two terms correspond to the sons' influence, the last can be controlled via SVD and provides us with the transfer matrices.

\mathcal{DH}^2 -matrix compression algorithm

Recursion following the cluster tree.

- Set up \mathcal{R}_{τ_C} during the downward pass.
- Construct V_{τ_C} , E_{τ'/τ_C} , and $V_{\tau_C}^* G_{\tau_C}$ during the upward pass.

Compared to \mathcal{H}^2 -matrices:

- Similar complexity estimates.
- Parallelization tricky, large clusters very time-consuming.
- Complexity $\mathcal{O}(n^2k)$

\mathcal{DH}^2 -matrix compression algorithm

Recursion following the cluster tree.

- Set up \mathcal{R}_{TC} during the downward pass.
- Construct V_{TC} , $E_{T'_{TC}}$, and $V_{TC}^* G_{TC}$ during the upward pass.

Compared to \mathcal{H}^2 -matrices:

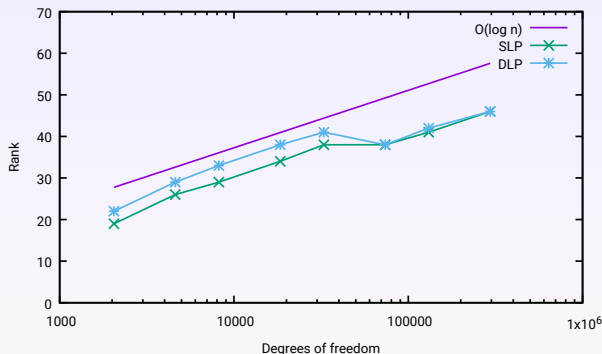
- Similar complexity estimates.
- Parallelization tricky, large clusters very time-consuming.
- Complexity $\mathcal{O}(n^2k)$, but
- efficient recompression of \mathcal{DH}^2 -matrices in $\mathcal{O}(nk^2 \log n)$.

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Experiment: Single and double layer potential

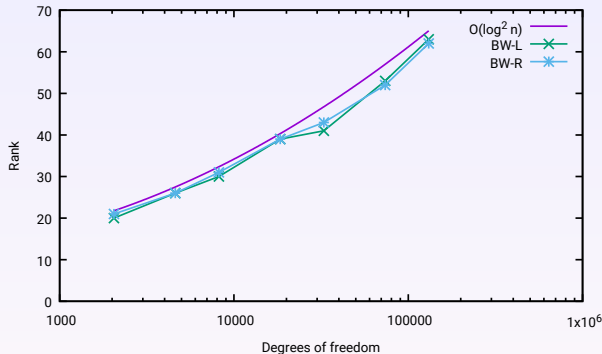
Approximate the matrix G corresponding to SLP and DLP operators for the unit sphere.



Result: Maximal rank grows like $\mathcal{O}(\log n)$.

Experiment: Preconditioner

Approximate the LR factorization of the Brakhage-Werner operator.



Result: Maximal rank grows like $\mathcal{O}(\log^2 n)$.

Conclusion

Challenge: Approximate high-frequency Helmholtz operators.

Approach: Directional approximation, \mathcal{DH}^2 -matrix.

→ Complexity $\mathcal{O}(nk^2 \log n)$.

Compression: Recursive algorithm, $\mathcal{O}(n^2k)$ operations.

Work in progress: On-the-fly recompression in $\mathcal{O}(nk^2 \log n)$,
 \mathcal{DH}^2 -matrix arithmetic operations, inclusion in H2Lib package

<http://www.h2lib.org>

