

# $\mathcal{DH}^2$ -matrix compression for Helmholtz problems

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# Overview

- 1 Introduction
- 2 Directional approximation
- 3 Algebraic compression
- 4 Recompression
- 5 Experiments

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# Helmholtz equation

**Goal:** Solve the Helmholtz equation

$$\begin{aligned}\Delta u(x) + \kappa^2 u(x) &= 0 && \text{for all } x \in \Omega, \\ u(x) &= f(x) && \text{for all } x \in \partial\Omega.\end{aligned}$$

**Challenge:** Indefinite operator, pollution effect.

**Approach:** Boundary integral formulation

$$u(x) = \int_{\partial\Omega} g(x, y) \frac{\partial u}{\partial n}(y) dy - \int_{\partial\Omega} \frac{\partial g}{\partial n(y)}(x, y) u(y) dy \quad \text{for all } x \in \Omega$$

with fundamental solution

$$g(x, y) = \frac{\exp(i\kappa\|x - y\|)}{4\pi\|x - y\|}.$$

**Challenge:** Standard discretization leads to dense matrix.

# Fast summation

Standard approach: Approximate  $g$  by degenerate kernel,

$$g(x, y) \approx \sum_{\nu=1}^k v_{\nu}(x) \sum_{\mu=1}^k s_{\nu\mu} \overline{w_{\mu}(y)}$$

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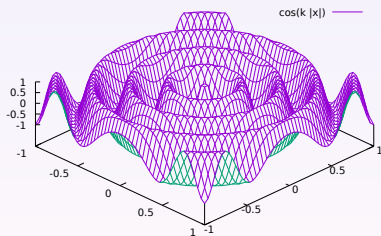
$$\int_{\partial\Omega} g(x, y) u(y) dy \approx \sum_{\nu=1}^k v_{\nu}(x) \int_{\partial\Omega} \sum_{\mu=1}^k s_{\nu\mu} \overline{w_{\mu}(y)} u(y) dy.$$

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Challenge: Helmholtz kernel oscillates rapidly if  $\kappa$  large.



$$g(x, y) = \frac{\exp(i\kappa\|x - y\|)}{\|x - y\|}.$$

Consequence: Standard approximation schemes require  $k \gg 1$ .

## Experiment: $\mathcal{H}^2$ -matrix compression

Helmholtz matrix approximated by quasi-optimal  $\mathcal{H}^2$ -matrix.

$n$	$\kappa$	$k$	$k/\kappa$	$M/n$
2048	8	33	4.1	12.5
4608	12	47	3.9	15.4
8192	16	61	3.8	18.3
18432	24	95	4.0	22.2
32768	32	132	4.1	25.4
73728	48	263	5.5	30.1
131072	64	333	5.2	36.0

**Observation:** Maximal rank appears proportional to wave number  $\kappa$ , storage looks like  $\mathcal{O}(n \log^2 \kappa)$ .



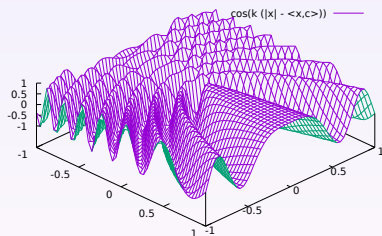
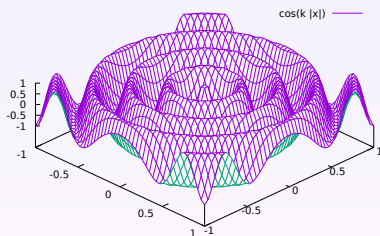
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# Directional approximation

Idea: Divide spherical wave by plane wave in direction  $c$ ,  $\|c\| = 1$ :

$$g(x, y) = \exp(i\kappa \langle x - y, c \rangle) \underbrace{\frac{\exp(i\kappa(\|x - y\| - \langle x - y, c \rangle))}{4\pi\|x - y\|}}_{=g_c(x, y)}.$$



Brandt (1991), Engquist/Ying (2007),  
Messner/Schanz/Darve (2012), Bebendorf/Kuske/Venn (2015)

# Directional interpolation

**Approach:** Apply interpolation to smoothed function  $g_c$ .

$$\tilde{g}_{\tau\sigma,c}(x, y) := \sum_{\nu,\mu=1}^k \mathcal{L}_{\tau,\nu}(x) g_c(\xi_{\tau,\nu}, \xi_{\sigma,\mu}) \mathcal{L}_{\sigma,\mu}(y).$$

**Directional approximation** in domain  $\tau \times \sigma$  given by

$$\begin{aligned} g(x, y) &= \exp(\iota\kappa\langle x - y, c \rangle) g_c(x, y) \approx \exp(\iota\kappa\langle x - y, c \rangle) \tilde{g}_{\tau\sigma,c}(x, y) \\ &= \sum_{\nu,\mu=1}^k \underbrace{\exp(\iota\kappa\langle x, c \rangle) \mathcal{L}_{\tau,\nu}(x)}_{=: \mathcal{L}_{\tau c,\nu}(x)} g_c(\xi_{\tau,\nu}, \xi_{\sigma,\mu}) \underbrace{\exp(-\iota\kappa\langle y, c \rangle) \mathcal{L}_{\sigma,\mu}(y)}_{=: \overline{\mathcal{L}_{\sigma c,\mu}(y)}} \\ &= \sum_{\nu,\mu=1}^k \mathcal{L}_{\tau c,\nu}(x) g_c(\xi_{\tau,\nu}, \xi_{\sigma,\mu}) \overline{\mathcal{L}_{\sigma c,\mu}(y)} =: \tilde{g}_{\tau\sigma}(x, y). \end{aligned}$$

# Matrix factorization

**Goal:** Approximate Galerkin matrix  $G \in \mathbb{C}^{\mathcal{I} \times \mathcal{I}}$  given by

$$g_{ij} = \int_{\partial\Omega} \varphi_i(x) \int_{\partial\Omega} g(x, y) \varphi_j(y) dy dx.$$

**Idea:** If  $g|_{\tau \times \sigma} \approx \tilde{g}_{\tau\sigma}$ , choose

$$\hat{\tau} := \{i \in \mathcal{I} : \text{supp } \varphi_i \subseteq \tau\}, \quad \hat{\sigma} := \{j \in \mathcal{I} : \text{supp } \varphi_j \subseteq \sigma\}$$

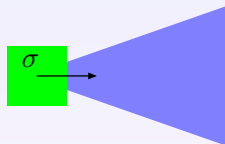
and approximate  $G|_{\hat{\tau} \times \hat{\sigma}}$  by discretizing  $\tilde{g}_{\tau\sigma}$  instead of  $g$ .

**Result:** Factorization  $G|_{\hat{\tau} \times \hat{\sigma}} \approx V_{\tau\mathcal{C}} S_{\tau\sigma} V_{\sigma\mathcal{C}}^*$  with

$$v_{\tau\mathcal{C}, i\nu} = \int_{\partial\Omega} \varphi_i(x) \mathcal{L}_{\tau\mathcal{C}, \nu}(x) dx, \quad s_{\tau\sigma, \nu\mu} = g_{\mathcal{C}}(\xi_{\tau, \nu}, \xi_{\sigma, \mu}).$$

# Directions

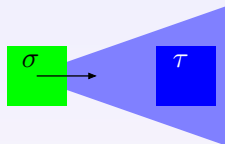
Kernel function  $g_c(\cdot, y)$  for  $y \in \sigma$   
is smooth in a cone with axis  $c$ .



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Interactions with clusters  $\tau$  inside  
the cone can be approximated.

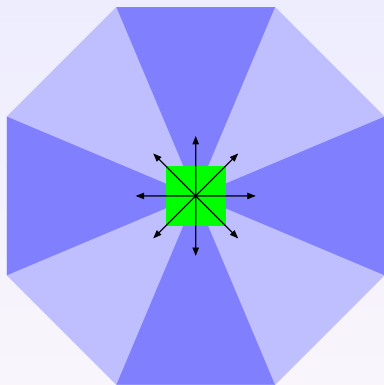


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Interactions with clusters  $\tau$  inside the cone can be approximated.

Multiple cones have to be used to cover the entire domain.

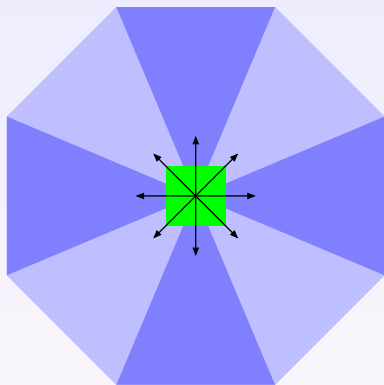


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Interactions with clusters  $\tau$  inside the cone can be approximated.

Multiple cones have to be used to cover the entire domain.



**Approach:** Choose fixed set of directions for each  $\sigma$ .

Analysis shows that  $\mathcal{O}(1 + \kappa^2 \text{diam}^2(\tau))$  directions are sufficient.



# Nested basis

**Problem:** Storing  $V_{\tau c}$  for all  $\tau, c$  requires  $\sim n^2$  units of memory.

**Solution:** Organize domains  $\tau$  in a **cluster tree** such that

$$\text{sons}(\tau) \neq \emptyset \quad \implies \quad \hat{\tau} = \bigcup_{\tau' \in \text{sons}(\tau)} \hat{\tau}',$$

and approximate Lagrange polynomials by weighted interpolation

$$\mathcal{L}_{\tau c, \nu} \approx \sum_{\nu'=1}^k \mathbf{e}_{\tau' \tau c, \nu' \nu} \mathcal{L}_{\tau' c', \nu'}.$$

**Result:** Nested representation  $V_{\tau c} |_{\hat{\tau}'} \approx V_{\tau' c'} E_{\tau' \tau c}$  if  $\tau' \in \text{sons}(\tau)$ .



Messner/Schanz/Darve (2012), B./Melenk (2015)

# $\mathcal{DH}^2$ -matrix

## Approach:

- Fix **set of directions**  $\mathcal{D}_\tau$  for all clusters  $\tau$ .
- Choose **admissible blocks**  $(\tau, \sigma, \mathbf{c}) \in \mathcal{L}_{I \times I}^+$ .
- Compute **leaf matrices**  $V_{\tau\mathbf{c}} \in \mathbb{C}^{\hat{\tau} \times k}$  for all leaves  $\tau$  and all  $\mathbf{c} \in \mathcal{D}_\tau$ .
- Find **transfer matrices**  $E_{\tau'\tau\mathbf{c}} \in \mathbb{C}^{k \times k}$  for all  $\tau' \in \text{sons}(\tau)$ ,  $\mathbf{c} \in \mathcal{D}_\tau$ .
- Create **coupling matrices**  $S_{\tau\sigma} \in \mathbb{C}^{k \times k}$  for all  $(\tau, \sigma, \mathbf{c}) \in \mathcal{L}_{I \times I}^+$ .

## Results:

- Storage  $\mathcal{O}(kn + k^2\kappa^2 \log n)$ .
- Matrix-vector multiplication in  $\mathcal{O}(kn + k^2\kappa^2 \log n)$  operations.
- High-frequency case:  $\kappa^2 \sim n$ , complexity  $\mathcal{O}(nk^2 \log n)$ .



B. (2015)

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# Approximation of general matrices

**Given:** Matrix  $G \in \mathbb{C}^{\mathcal{I} \times \mathcal{I}}$  and cluster/block structure.

**Goal:** Find  $\mathcal{DH}^2$ -matrix approximation of  $G$ ,  
i.e., directional cluster basis ( $V_{\tau\mathcal{C}}$ ) and coupling matrices ( $S_{\tau\sigma}$ ).

**Orthogonal basis:** If we ensure  $V_{\tau\mathcal{C}}^* V_{\tau\mathcal{C}} = I$ , the best coupling matrices are given by orthogonal projection

$$G|_{\hat{\tau} \times \hat{\sigma}} \approx V_{\tau\mathcal{C}} \underbrace{V_{\tau\mathcal{C}}^* G|_{\hat{\tau} \times \hat{\sigma}} V_{\sigma\mathcal{C}}}_{=: S_{\tau\sigma}} V_{\sigma\mathcal{C}}^*$$

Row and column projections can be analyzed separately,

$$\begin{aligned} \|G|_{\hat{\tau} \times \hat{\sigma}} - V_{\tau\mathcal{C}} S_{\tau\sigma} V_{\sigma\mathcal{C}}^*\|_2^2 &\leq \|G|_{\hat{\tau} \times \hat{\sigma}} - V_{\tau\mathcal{C}} V_{\tau\mathcal{C}}^* G|_{\hat{\tau} \times \hat{\sigma}}\|_2^2 \\ &\quad + \|G|_{\hat{\tau} \times \hat{\sigma}} - G|_{\hat{\tau} \times \hat{\sigma}} V_{\sigma\mathcal{C}} V_{\sigma\mathcal{C}}^*\|_2^2. \end{aligned}$$

# Low-rank structure

**Question:** What has to be approximated by  $V_{\tau\mathbf{c}}$ ?

**Answer:** All blocks  $G|_{\hat{\tau} \times \hat{\sigma}}$  with  $(\tau, \sigma, \mathbf{c}) \in \mathcal{L}_{I \times I}^+$ .

# Low-rank structure

**Question:** What has to be approximated by  $V_{\tau c}$ ?

**Partial answer:** All blocks  $G|_{\hat{\tau} \times \hat{\sigma}}$  with  $(\tau, \sigma, c) \in \mathcal{L}_{I \times I}^+$ .

**Important:** Directional cluster basis is **nested**,  $V_{\tau c}|_{\hat{\tau}'} = V_{\tau' c'} E_{\tau' \tau c}$ .  
Anything that is not approximated in the son clusters cannot be approximated in the father.

# Low-rank structure

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Anything that is not approximated in the son clusters cannot be approximated in the father.

**Complete answer:** All blocks  $G|_{\hat{\tau} \times \hat{\sigma}^*}$  where

- $(\tau^*, \sigma^*, c^*) \in \mathcal{L}_{I \times I}^+$ ,
- $\tau$  is a descendant of  $\tau^*$ , and
- $c$  approximates  $c^*$ .

We collect these tuples in sets  $\mathcal{R}_{\tau c} \subseteq \mathcal{L}_{I \times I}^+$ .

# Leaf matrices

Assume that  $\tau$  is a leaf cluster and that  $\mathcal{R}_{\tau c}$  is available for all  $c \in \mathcal{D}_\tau$ .

**Goal:** Find orthogonal  $V_{\tau c}$  with  $k$  columns such that

$$G|_{\hat{\tau} \times \hat{\sigma}^*} \approx V_{\tau c} V_{\tau c}^* G|_{\hat{\tau} \times \hat{\sigma}^*} \quad \text{for all } (\tau^*, \sigma^*, c^*) \in \mathcal{R}_{\tau c}.$$



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**Reformulation:** Define the submatrix

$$G_{\tau c} := G|_{\hat{\tau} \times \mathcal{F}_{\tau c}}, \quad \mathcal{F}_{\tau c} := \bigcup \{ \hat{\sigma}^* : (\tau^*, \sigma^*, c^*) \in \mathcal{R}_{\tau c} \}$$

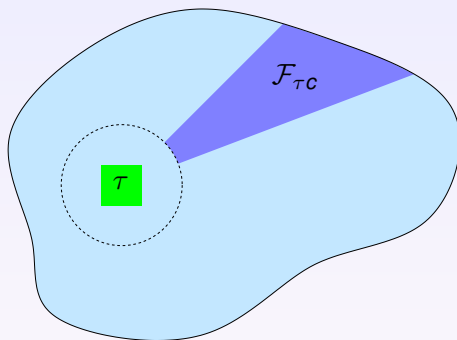
with **farfield matrix**  $G_{\tau c}$  and consider only

$$G_{\tau c} \approx V_{\tau c} V_{\tau c}^* G_{\tau c}.$$

→ Compute SVD of  $G_{\tau c}$ , remove small singular values.

# Directional farfield

Question: What does  $\mathcal{F}_{\tau c}$  represent?



Answer:  $\mathcal{F}_{\tau c}$  corresponds to basis functions whose supports are

- sufficiently far from  $\tau$  (i.e.,  $\text{dist} \geq \text{diam}(\tau), \kappa \text{diam}(\tau)^2$ ), and
- contained in a cone around the direction  $c$ .

# Transfer matrices

Assume  $\text{sons}(\tau) = \{\tau_1, \tau_2\}$ ,  $\mathbf{c}_1 \in \mathcal{D}_{\tau_1}$ ,  $\mathbf{c}_2 \in \mathcal{D}_{\tau_2}$  approximate  $\mathbf{c} \in \mathcal{D}_{\tau}$ .

**Recursion:** Ensure that  $V_{\tau_1 \mathbf{c}_1}$ ,  $V_{\tau_2 \mathbf{c}_2}$  have already been computed.

**Reduced matrices** given by

$$\widehat{\mathbf{G}}_{\tau \mathbf{c}} = \begin{pmatrix} \mathbf{V}_{\tau_1 \mathbf{c}_1}^* \mathbf{G}|_{\hat{\tau}_1 \times \mathcal{F}_{\tau \mathbf{c}}} \\ \mathbf{V}_{\tau_2 \mathbf{c}_2}^* \mathbf{G}|_{\hat{\tau}_2 \times \mathcal{F}_{\tau \mathbf{c}}} \end{pmatrix}, \quad \widehat{\mathbf{V}}_{\tau \mathbf{c}} = \begin{pmatrix} \mathbf{E}_{\tau_1 \tau \mathbf{c}} \\ \mathbf{E}_{\tau_2 \tau \mathbf{c}} \end{pmatrix}$$

lead to

$$\begin{aligned} \|\mathbf{G}_{\tau \mathbf{c}} - \mathbf{V}_{\tau \mathbf{c}} \mathbf{V}_{\tau \mathbf{c}}^* \mathbf{G}_{\tau \mathbf{c}}\|_2^2 &\leq \|\mathbf{G}|_{\hat{\tau}_1 \times \mathcal{F}_{\tau \mathbf{c}}} - \mathbf{V}_{\tau_1 \mathbf{c}_1} \mathbf{V}_{\tau_1 \mathbf{c}_1}^* \mathbf{G}|_{\hat{\tau}_1 \times \mathcal{F}_{\tau \mathbf{c}}}\|_2^2 \\ &\quad + \|\mathbf{G}|_{\hat{\tau}_2 \times \mathcal{F}_{\tau \mathbf{c}}} - \mathbf{V}_{\tau_2 \mathbf{c}_2} \mathbf{V}_{\tau_2 \mathbf{c}_2}^* \mathbf{G}|_{\hat{\tau}_2 \times \mathcal{F}_{\tau \mathbf{c}}}\|_2^2 \\ &\quad + \|\widehat{\mathbf{G}}_{\tau \mathbf{c}} - \widehat{\mathbf{V}}_{\tau \mathbf{c}} \widehat{\mathbf{V}}_{\tau \mathbf{c}}^* \widehat{\mathbf{G}}_{\tau \mathbf{c}}\|_2^2. \end{aligned}$$

The first two terms correspond to the sons' influence, the last can be controlled via SVD and provides us with the transfer matrices.

# $\mathcal{DH}^2$ -matrix compression algorithm

Recursion following the cluster tree.

- Set up  $\mathcal{R}_{\tau_C}$  during the downward pass.
- Construct  $V_{\tau_C}$ ,  $E_{\tau'/\tau_C}$ , and  $V_{\tau_C}^* G_{\tau_C}$  during the upward pass.

Compared to  $\mathcal{H}^2$ -matrices:

- Similar complexity estimates.
- Loss of sparsity, some error estimates no longer work.
- Parallelization tricky, large clusters very time-consuming.
- Complexity  $\mathcal{O}(n^2 k)$ .

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# Approximation of $\mathcal{DH}^2$ -matrices

**Challenge:** Compression of a general matrix takes  $\mathcal{O}(n^2k)$  operations.

**Idea:** Obtain intermediate  $\mathcal{DH}^2$ -matrix by interpolation and **recompress** to reduce the rank.

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**Nested structure:** We have

$$(V_{\tau c} S_{\tau\sigma} V_{\sigma c}^*)|_{\hat{\tau}'} = V_{\tau' c'} E_{\tau' \tau c} S_{\tau\sigma} V_{\sigma c}^*,$$

and induction yields

$$G|_{\hat{\tau} \times \hat{\sigma}^*} \approx V_{\tau c} E_{\tau \tau^* c^*} S_{\tau^* \sigma^*} V_{\sigma^* c^*}^* \quad \text{for all } (\tau^*, \sigma^*, c^*) \in \mathcal{R}_{\tau c}$$

with suitable “long-range” transfer matrices  $E_{\tau \tau^* c^*}$ .

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**Result:** Intermediate low-rank approximation  $G_{\tau c} \approx V_{\tau c} B_{\tau c}^*$ .



# Weighted truncation

**Goal:** Find optimized low-rank factorization of intermediate approximation

$$G_{TC} \approx V_{TC} B_{TC}^*.$$

**Idea:** Since approximation depends only on singular values and left singular vectors, use skinny Householder factorization

$$B_{TC} = P_{TC} Z_{TC}$$

with weight matrix  $Z_{TC} \in \mathbb{C}^{k \times k}$  and work with

$$G_{TC} P_{TC}^* \approx V_{TC} Z_{TC}^*.$$

**Result:** If all weights  $(Z_{TC})_{T,C}$  are available, compression requires only  $\mathcal{O}(nk^2 \log n)$  operations instead of  $\sim n^2 k$ .

# Weight matrices

**Goal:** Compute weight matrices with  $B_{\tau c} = P_{\tau c} Z_{\tau c}$ .

**Idea:** Take advantage of nested structure,  $\tau \in \text{sons}(\tau^+)$ .

$$B_{\tau c} = \begin{pmatrix} V_{\sigma_1 c} S_{\tau \sigma_1}^* \\ \vdots \\ V_{\sigma_s c} S_{\tau \sigma_s}^* \\ B_{\tau^+ c_1^+} E_{\tau \tau^+ c_1^+}^* \\ \vdots \\ B_{\tau^+ c_\ell^+} E_{\tau \tau^+ c_\ell^+}^* \end{pmatrix}$$

with blocks  $(\tau, \sigma_1, c), \dots, (\tau, \sigma_s, c) \in \mathcal{L}_{\mathcal{I} \times \mathcal{I}}^+$ , and the father's directions  $c_1^+, \dots, c_\ell^+ \in \mathcal{D}_{\tau^+}$  approximating  $c$ .

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**Result:** Construction of all  $(Z_{\tau c})_{\tau, c}$  in  $\mathcal{O}(nk^2 \log n)$  operations.

# Recompression algorithm

**Directional interpolation** yields intermediate  $\mathcal{DH}^2$ -matrix approximation that is used **implicitly** and not stored completely.

**Farfield matrices**  $G_{TC}$  are replaced by low-rank factorizations  $V_{TC}Z_{TC}^*$  with suitable weight matrices.

**Weight matrices** constructed by top-down recursion and Householder factorization.

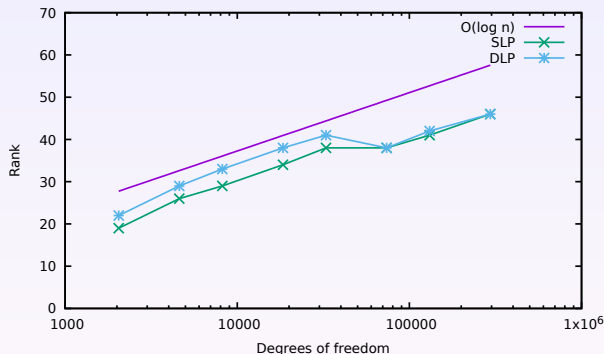
**Result:** Compression in  $\mathcal{O}(nk^2 \log n)$  operations.

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# Experiment: Single and double layer potential

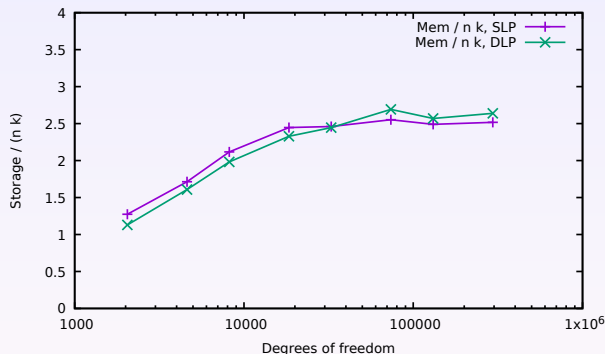
Approximate the matrix  $G$  corresponding to SLP and DLP operators for the unit sphere.



**Result:** Maximal rank grows like  $O(\log n)$ .

# Experiment: Single and double layer potential

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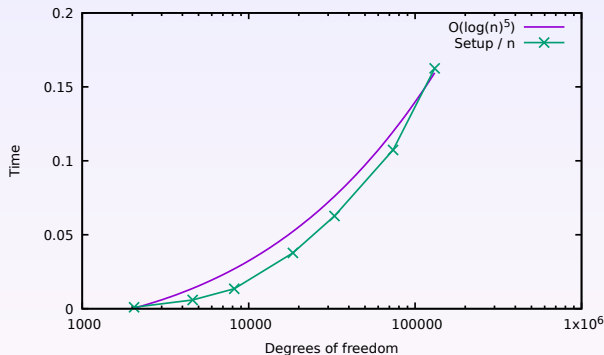


**Result:** Storage grows like  $\mathcal{O}(nk)$ .



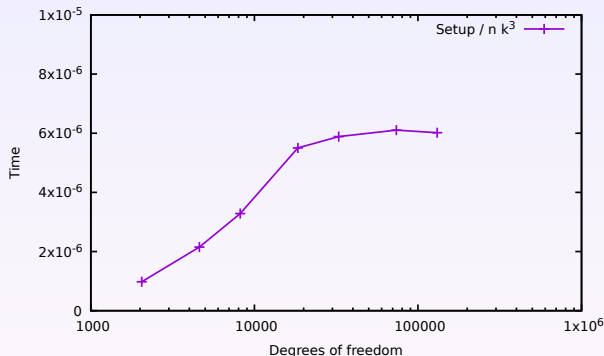
# Experiment: Recompression

Approximate the matrix  $G$  corresponding to the SLP operator by directional interpolation, followed by algebraic recompression.



# Experiment: Recompression

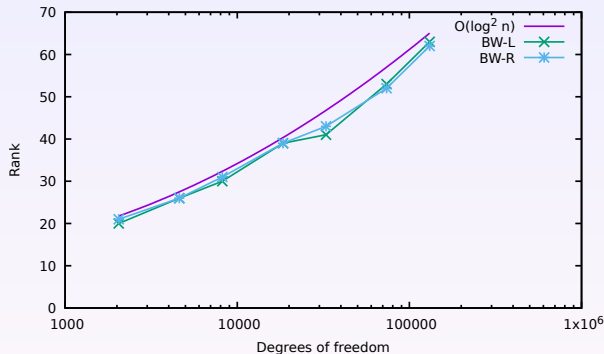
Approximate the matrix  $G$  corresponding to the SLP operator by directional interpolation, followed by algebraic recompression.



**Result:** Time appears to grow like  $\mathcal{O}(nk^3)$ .

# Experiment: Preconditioner

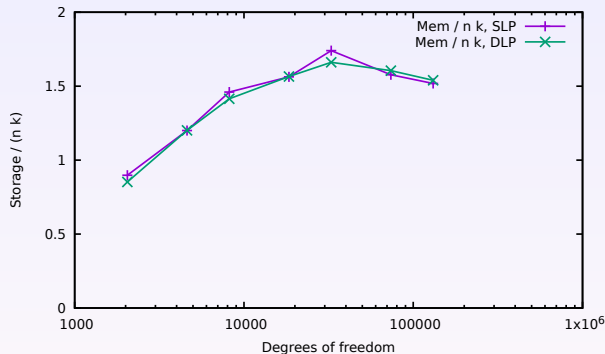
Approximate the LR factorization of the Brakhage-Werner operator.



**Result:** Maximal rank grows like  $\mathcal{O}(\log^2 n)$ .

# Experiment: Preconditioner

Approximate the LR factorization of the Brakhage-Werner operator.



**Result:** Storage grows like  $\mathcal{O}(nk)$ .

# Conclusion

**Challenge:** Approximate high-frequency Helmholtz operators.

**Approach:** Directional approximation leads to  $\mathcal{DH}^2$ -matrix.

→ Complexity  $\mathcal{O}(nk^2 \log n)$ .

**Compression:** Recursive algorithm,  $\mathcal{O}(n^2k)$  operations.

**Recompression:** Combine directional interpolation with on-the-fly compression,  $\mathcal{O}(nk^2 \log n)$  operations.

**Work in progress:** Arithmetic operations, inclusion in H2Lib package

<http://www.h2lib.org>

